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Transformer-based Algorithmics

**Malte Luttermann – Institute for Humanities-Centered
Artificial Intelligence (CHAI)**

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Topics for Today

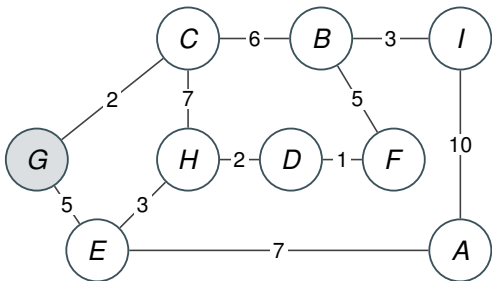
- Dijkstra's algorithm
- A* algorithm

Acknowledgement

- Some of the upcoming content is taken from the following lectures and corresponding exercises and has been translated and partially modified:
 - Prof. Dr. Ralf Möller: »[Algorithmen und Datenstrukturen – Vorlesung](#)«
 - Dr. Magnus Bender: »[Algorithmen und Datenstrukturen – Übung](#)«

Exercise 1 – Dijkstra's Algorithm

Consider the weighted (undirected) graph $G = (V, E)$ below.



Apply **Dijkstra's algorithm** from vertex G . After each while-loop iteration, state:

- the extracted vertex u (with $u.d$ and $u.parent$),
- the contents of the priority queue q (each v with $v.d$, $v.parent$).

Tie-breaking: on equal distances, keep the original order.

Finally, give the **shortest path** $G \rightarrow I$.

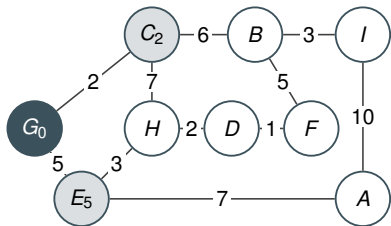
Dijkstra's Algorithm – Reminder

Dijkstra in a Nutshell

Computes **single-source shortest paths** in a graph with *non-negative* edge weights.

- (1) Initialize $s.d = 0$, $s.parent = s$, and $v.d = \infty$ for every other vertex. Insert s into a priority queue q (key = tentative distance).
 - (2) While q is non-empty: extract the vertex u with smallest $u.d$ and *relax* every edge $u \rightarrow v$: if $u.d + w(u, v) < v.d$, set $v.d = u.d + w(u, v)$ and $v.parent = u$ (insert/update v in q).
- Once u is extracted from q , its distance $u.d$ is **final** – it will never decrease.
 - The parent pointers form the **shortest-path tree** rooted at s .
 - Runtime with a binary heap: $O((|V| + |E|) \log |V|)$.

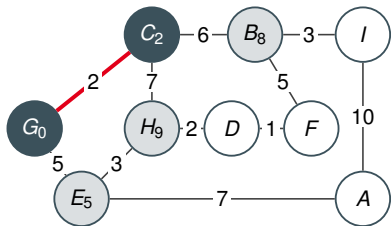
Exercise 1 – Dijkstra Trace from G



Step 1. Extract source G ; discover C, E .

Extracted from queue			Current queue q		
u	$u.d$	$u.parent$	v	$v.d$	$v.parent$
G	0	G	C	2	G
			E	5	G

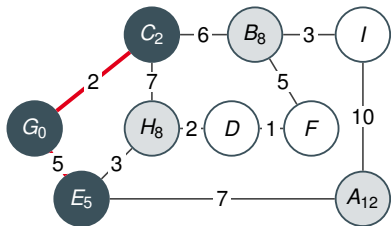
Exercise 1 – Dijkstra Trace from G



Step 2. Extract C ; discover B, H .

Extracted from queue			Current queue q		
u	$u.d$	$u.parent$	v	$v.d$	$v.parent$
G	0	G	C	2	G
			E	5	G
C	2	G	E	5	G
			B	8	C
			H	9	C

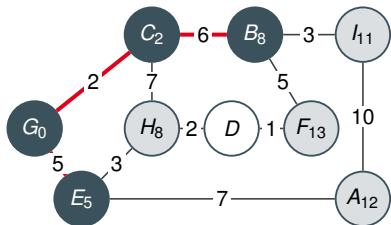
Exercise 1 – Dijkstra Trace from G



Step 3. Extract E ; **relax** H : $9 \rightarrow 8$ (parent E); discover A .

Extracted from queue			Current queue q		
u	$u.d$	$u.parent$	v	$v.d$	$v.parent$
G	0	G	C	2	G
			E	5	G
C	2	G	E	5	G
			B	8	C
			H	9	C
E	5	G	B	8	C
			H	8	E
			A	12	E

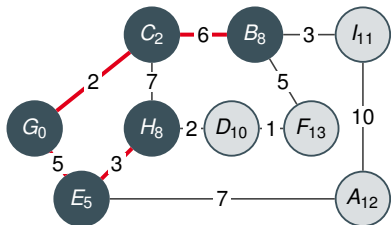
Exercise 1 – Dijkstra Trace from G



Step 4. Extract B ; discover I, F .

Extracted from queue			Current queue q		
u	$u.d$	$u.parent$	v	$v.d$	$v.parent$
G	0	G	C	2	G
			E	5	G
C	2	G	E	5	G
			B	8	C
			H	9	C
E	5	G	B	8	C
			H	8	E
			A	12	E
B	8	C	H	8	E
			I	11	B
			A	12	E
			F	13	B

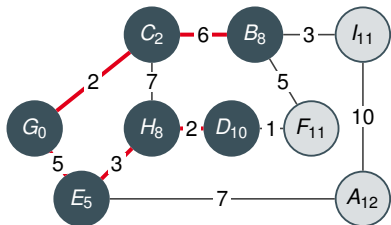
Exercise 1 – Dijkstra Trace from G



Step 5. Extract H ; discover D .

Extracted from queue			Current queue q		
u	$u.d$	$u.parent$	v	$v.d$	$v.parent$
G	0	G	C	2	G
			E	5	G
C	2	G	E	5	G
			B	8	C
			H	9	C
E	5	G	B	8	C
			H	8	E
			A	12	E
B	8	C	H	8	E
			I	11	B
			A	12	E
			F	13	B
H	8	E	D	10	H
			I	11	B
			A	12	E
			F	13	B

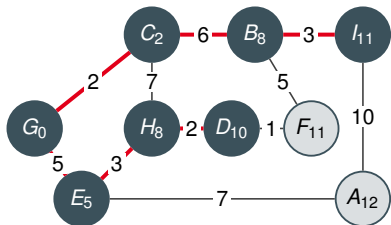
Exercise 1 – Dijkstra Trace from G



Step 6. Extract D ; **relax** F : $13 \rightarrow 11$
(parent D).

Extracted from queue			Current queue q		
u	$u.d$	$u.parent$	v	$v.d$	$v.parent$
G	0	G	C	2	G
			E	5	G
C	2	G	E	5	G
			B	8	C
			H	9	C
E	5	G	B	8	C
			H	8	E
			A	12	E
B	8	C	H	8	E
			I	11	B
			A	12	E
			F	13	B
H	8	E	D	10	H
			I	11	B
			A	12	E
			F	13	B
D	10	H	I	11	B
			F	11	D
			A	12	E

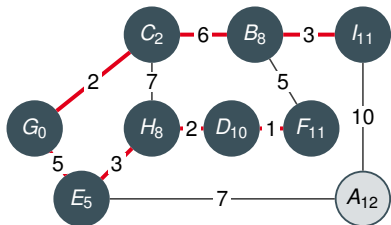
Exercise 1 – Dijkstra Trace from G



Step 7. Extract I ; no neighbors improved.

Extracted from queue			Current queue q		
u	$u.d$	$u.parent$	v	$v.d$	$v.parent$
G	0	G	C	2	G
			E	5	G
C	2	G	E	5	G
			B	8	C
			H	9	C
E	5	G	B	8	C
			H	8	E
			A	12	E
B	8	C	H	8	E
			I	11	B
			A	12	E
			F	13	B
H	8	E	D	10	H
			I	11	B
			A	12	E
			F	13	B
D	10	H	I	11	B
			F	11	D
			A	12	E
I	11	B	F	11	D
			A	12	E

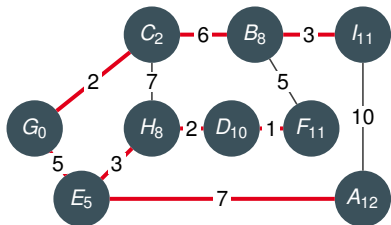
Exercise 1 – Dijkstra Trace from G



Step 8. Extract F ; no neighbors improved.

Extracted from queue			Current queue q		
u	$u.d$	$u.parent$	v	$v.d$	$v.parent$
G	0	G	C	2	G
			E	5	G
C	2	G	E	5	G
			B	8	C
			H	9	C
E	5	G	B	8	C
			H	8	E
			A	12	E
B	8	C	H	8	E
			I	11	B
			A	12	E
			F	13	B
H	8	E	D	10	H
			I	11	B
			A	12	E
			F	13	B
D	10	H	I	11	B
			F	11	D
			A	12	E
I	11	B	F	11	D
			A	12	E
F	11	D	A	12	E

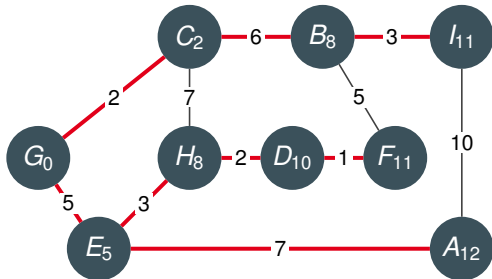
Exercise 1 – Dijkstra Trace from G



Step 9. Extract A. Queue empty – **done**.

Extracted from queue			Current queue q		
u	$u.d$	$u.parent$	v	$v.d$	$v.parent$
G	0	G	C	2	G
			E	5	G
C	2	G	E	5	G
			B	8	C
			H	9	C
E	5	G	B	8	C
			H	8	E
			A	12	E
B	8	C	H	8	E
			I	11	B
			A	12	E
			F	13	B
H	8	E	D	10	H
			I	11	B
			A	12	E
			F	13	B
D	10	H	I	11	B
			F	11	D
			A	12	E
I	11	B	F	11	D
			A	12	E
F	11	D	A	12	E
A	12	E	(empty)		

Exercise 1 – Result: Shortest-Path Tree from G



Vertex labels v_d show the final distance from G ; **red** edges form the **shortest-path tree**.

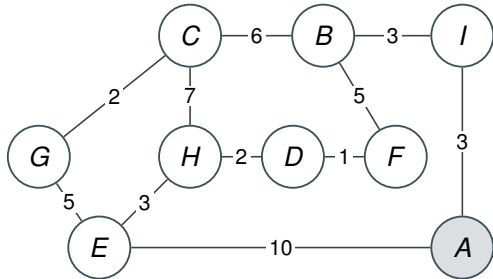
Shortest path $G \rightarrow I$:

$$G \xrightarrow{2} C \xrightarrow{6} B \xrightarrow{3} I$$

Total length: 11

Exercise 2 – A* Algorithm

Find the shortest path $A \rightarrow G$ in the (undirected) graph below using A* with the heuristic h_1 .



Heuristic h_1 (estimate to G):

k	A	B	C	D	E	F	G	H	I
$h_1(k)$	6	6	1	7	2	6	0	1	11

After each iteration, give $k.g$, $k.f$ and $k.parent$ for each vertex.

Tie-breaking: on equal f , pick the lexicographical smallest vertex.

A* Algorithm – Reminder

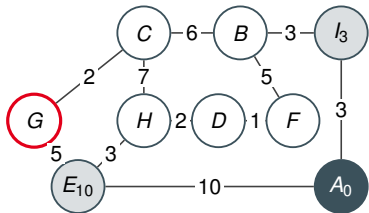
A* in a Nutshell

Best-first search guided by the **evaluation function** $f(n) = g(n) + h(n)$, where $g(n)$ is the cost from start s to n and $h(n)$ estimates the remaining cost from n to goal G . Maintain an open priority queue (key = f). Repeat: extract u with smallest $u.f$; if $u = G$, stop. Otherwise relax every edge $u \rightarrow v$: if $u.g + w(u, v) < v.g$, set $v.g, v.f = v.g + h(v)$, $v.parent = u$.

Heuristic properties

- **Admissible**: $h(n) \leq d^*(n, G)$ for all $n \Rightarrow A^*$ finds an optimal path.
- **Optimal**: $h(n) = d^*(n, G)$ – the true remaining distance (perfectly informed).
- $h \equiv 0$: no guidance $\Rightarrow A^*$ degenerates to **Dijkstra**.

Exercise 2 – A^* Trace from A to G



Step 1. Extract A ($f=6$); discover E, I .

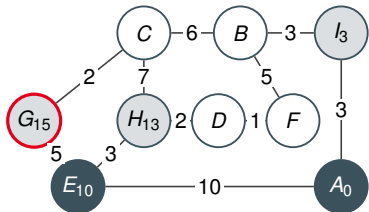
k	$h(k)$	$k.g$	$k.f$	$k.parent$
A	6	0	6	—
B	6			
C	1			
D	7			
E	2	10	12	A
F	6			
G	0			
H	1			
I	11	3	14	A

current

in open queue

closed

Exercise 2 – A* Trace from A to G



Step 2. Extract E ($f=12$); discover G , H .

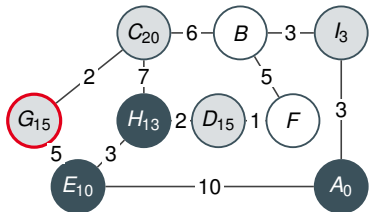
k	$h(k)$	$k.g$	$k.f$	$k.parent$
A	6	0	6	—
B	6			
C	1			
D	7			
E	2	10	12	A
F	6			
G	0	15	15	E
H	1	13	14	E
I	11	3	14	A

current

in open queue

closed

Exercise 2 – A* Trace from A to G



Step 3. Extract H ($f=14$, lex. before I);
discover C, D .

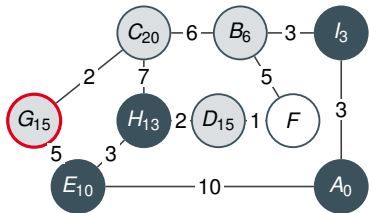
k	$h(k)$	$k.g$	$k.f$	$k.parent$
A	6	0	6	—
B	6			
C	1	20	21	H
D	7	15	22	H
E	2	10	12	A
F	6			
G	0	15	15	E
H	1	13	14	E
I	11	3	14	A

current

in open queue

closed

Exercise 2 – A^* Trace from A to G



Step 4. Extract I ($f=14$); discover B .

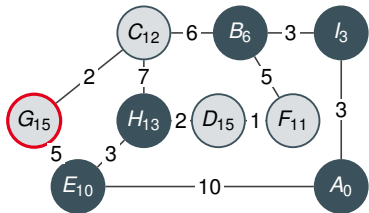
k	$h(k)$	$k.g$	$k.f$	$k.parent$
A	6	0	6	—
B	6	6	12	I
C	1	20	21	H
D	7	15	22	H
E	2	10	12	A
F	6			
G	0	15	15	E
H	1	13	14	E
I	11	3	14	A

current

in open queue

closed

Exercise 2 – A* Trace from A to G



Step 5. Extract B ($f=12$); **update** C
($21 \rightarrow 13$); discover F .

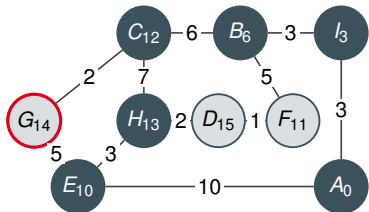
k	$h(k)$	$k.g$	$k.f$	$k.parent$
A	6	0	6	—
B	6	6	12	I
C	1	12	13	B
D	7	15	22	H
E	2	10	12	A
F	6	11	17	B
G	0	15	15	E
H	1	13	14	E
I	11	3	14	A

current

in open queue

closed

Exercise 2 – A* Trace from A to G



Step 6. Extract C ($f=13$); **update** G
($15 \rightarrow 14$).

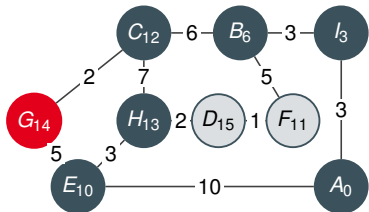
k	$h(k)$	$k.g$	$k.f$	$k.parent$
A	6	0	6	—
B	6	6	12	I
C	1	12	13	B
D	7	15	22	H
E	2	10	12	A
F	6	11	17	B
G	0	14	14	C
H	1	13	14	E
I	11	3	14	A

current

in open queue

closed

Exercise 2 – A^* Trace from A to G



Step 7. Extract G ($f=14$) – **goal reached!**

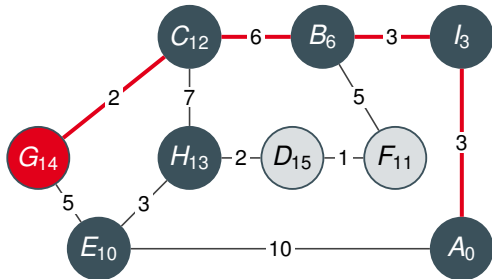
k	$h(k)$	$k.g$	$k.f$	$k.parent$
A	6	0	6	—
B	6	6	12	I
C	1	12	13	B
D	7	15	22	H
E	2	10	12	A
F	6	11	17	B
G	0	14	14	C
H	1	13	14	E
I	11	3	14	A

current

in open queue

closed

Exercise 2 – Result: Shortest Path $A \rightarrow G$



Vertex labels v_g show the final g -value.

Shortest path $A \rightarrow G$:

$$A \xrightarrow{3} I \xrightarrow{3} B \xrightarrow{6} C \xrightarrow{2} G$$

Total length: $3 + 3 + 6 + 2 = 14$

Nodes D and F were *discovered* but never expanded – A^* pruned them because their f -values exceeded the cost of the goal.

Exercise 3 – Evaluating Four Heuristics

Which of the following heuristics are admissible? Which are optimal?

k	A	B	C	D	E	F	G	H	I
$h_1(k)$	6	6	1	7	2	6	0	1	11
$h_2(k)$	12	8	2	10	5	11	0	8	15
$h_3(k)$	12	8	2	10	5	11	0	8	11
$h_4(k)$	0	0	0	0	0	0	0	0	0
$d^*(k, G)$	14	8	2	10	5	11	0	8	11

Exercise 3 – Evaluating Four Heuristics

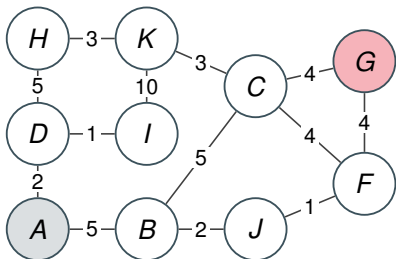
Which of the following heuristics are admissible? Which are optimal?

k	A	B	C	D	E	F	G	H	I
$h_1(k)$	6	6	1	7	2	6	0	1	11
$h_2(k)$	12	8	2	10	5	11	0	8	15
$h_3(k)$	12	8	2	10	5	11	0	8	11
$h_4(k)$	0	0	0	0	0	0	0	0	0
$d^*(k, G)$	14	8	2	10	5	11	0	8	11

- h_1 : admissible (all $\leq d^*$), not optimal.
- h_2 : **not admissible** – overestimates I ($15 > 11$); A^* would commit to the wrong path ($f(G)=15 < 18 = f(I)$ after A, E expanded).
- h_3 : matches d^* except at A ($12 < 14$); admissible. Setting $h_3(A)=14$ makes it optimal.
- $h_4 \equiv 0$: trivially admissible – exactly Dijkstra.

Exercise 4 – Heuristic Design for A*

Consider a different weighted graph (10 vertices, target G).



Given heuristic h (estimate to G):

k	A	B	C	D	F	G	H	I	J	K
$h(k)$	16	5	1	15	4	0	6	11	5	4

- Assess h regarding **admissibility** and **optimality**.
- Give an admissible *and* an optimal heuristic.
- How are A* and Dijkstra related?

Exercise 4 – Solution

(a) The given h is **neither admissible nor optimal**.

- Not admissible: $h(A)=16 > 12 = d^*(A, G)$ and $h(D)=15 > 14 = d^*(D, G)$ – *overestimates*.
- Not optimal: $h(C)=1 < 4 = d^*(C, G)$ – *underestimates*.

(b) Optimal and admissible heuristics:

k	A	B	C	D	F	G	H	I	J	K
$h_{\text{opt}}=d^*$	12	7	4	14	4	0	10	15	5	7
h_{adm}	12	5	1	14	4	0	6	11	5	4

Any h with $0 \leq h(k) \leq d^*(k, G)$ is admissible.

(c) **A* vs. Dijkstra.** Dijkstra is exactly A* with $h \equiv 0$ – it expands purely by g . An informed admissible heuristic guides A* towards the goal while preserving optimality.