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Transformer-based Algorithmics

**Malte Luttermann – Institute for Humanities-Centered
Artificial Intelligence (CHAI)**

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Topics for Today

- Hashing with open addressing
 - Linear probing
 - Quadratic probing

Acknowledgement

- Some of the upcoming content is taken from the following lectures and corresponding exercises and has been translated and partially modified:
 - Prof. Dr. Ralf Möller: »[Algorithmen und Datenstrukturen – Vorlesung](#)«
 - Dr. Magnus Bender: »[Algorithmen und Datenstrukturen – Übung](#)«

Exercise 1 – Hashing

Insert the keys (in this order):

40	12	7	26	34	51	33
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into an initially empty hash table of length 9, using $h(k) = k \bmod 9$. Use **internal collision handling** (open addressing) with:

- (a) **Linear probing**: $h_i(k) = (h(k) + i) \bmod 9$
- (b) **Quadratic probing**: $h_i(k) = (h(k) + i^2) \bmod 9$

Open Addressing – Reminder

Probe step $i = 0, 1, 2, \dots$ – try slot $h_i(k)$. On collision (slot occupied), increment i and retry until an empty slot is found.

Report the **final table** and all collisions (**X**) along the way.

Exercise 1 (a) – Linear Probing: First Insertions

Step	0	1	2	3	4	5	6	7	8
$h(40) = 4$					40				
$h(12) = 3$				12	40				
$h(7) = 7$				12	40			7	
$h(26) = 8$				12	40			7	26
$h(34) = 7$				12	40			X	26
$h_1(34) = 8$				12	40			7	X
$h_2(34) = 0$	34			12	40			7	26

Inserting 34: slot 7 holds 7 (X), slot 8 holds 26 (X), wrap to slot 0 – empty $\Rightarrow$
 place 34 at index 0.

Exercise 1 (a) – Linear Probing: Final Insertions & Result

Step	0	1	2	3	4	5	6	7	8
$h(51) = 6$	34			12	40		51	7	26
$h(33) = 6$	34			12	40		X	7	26
$h_1(33) = 7$	34			12	40		51	X	26
$h_2(33) = 8$	34			12	40		51	7	X
$h_3(33) = 0$	X			12	40		51	7	26
$h_4(33) = 1$	34	33		12	40		51	7	26

Final table:

34	33		12	40		51	7	26
0	1	2	3	4	5	6	7	8

Inserting 33 caused **4 collisions** – primary clustering at slots 6–8, 0.

Exercise 1 (b) – Quadratic Probing

Step	0	1	2	3	4	5	6	7	8
$h(40) = 4$					40				
$h(12) = 3$				12	40				
$h(7) = 7$				12	40			7	
$h(26) = 8$				12	40			7	26
$h(34) = 7$				12	40			X	26
$h_1(34) = 8$				12	40			7	X
$h_2(34) = 2$			34	12	40			7	26
$h(51) = 6$			34	12	40		51	7	26
$h(33) = 6$			34	12	40		X	7	26
$h_1(33) = 7$			34	12	40		51	X	26
$h_2(33) = 1$		33	34	12	40		51	7	26

Final:

	33	34	12	40		51	7	26
0	1	2	3	4	5	6	7	8

Exercise 1 – Comparison & Discussion

(a) Linear Probing

34	33		12	40		51	7	26
0	1	2	3	4	5	6	7	8

Collisions: 34: 2, 33: 4 (6 total)

(b) Quadratic Probing

	33	34	12	40		51	7	26
0	1	2	3	4	5	6	7	8

Collisions: 34: 2, 33: 2 (4 total)

Observations

- **Primary clustering** (linear): the run 6, 7, 8, 0 extends with every key hashing into it – 33 collides 4 times.
- Quadratic offsets 1, 4, 9, ... break primary clustering and halve the cascade.
- *Caveat:* for $m = 9$, $i^2 \bmod 9 \in \{0, 1, 4, 7\}$ – not all slots reachable. $\Rightarrow$ choose **m prime**.
- Both share **secondary clustering**: same $h(k) \Rightarrow$ same probe path.

Exercise 2 – Designing a Hash Function

For three-letter month names, the following hash function is proposed:

$$h(u) = (\text{pos}(u_1) + \text{pos}(u_2)) \bmod m,$$

where $\text{pos}(u_i)$ is the position of the i -th letter of u in the alphabet (i.e., $a = 1, b = 2, \dots, z = 26$).

Month names: jan, feb, mar, apr, mai, jun, jul.

Question. Can collisions occur with this hash function? If yes, give an example.

Exercise 2 – Solution & Discussion

Yes – collisions are unavoidable. The function uses only the first *two* letters, so any two months sharing those letters collide **for every** m :

$$h(\text{jun}) = h(\text{jul}) = (10 + 21) \bmod m = 31 \bmod m.$$

Pre-mod values for all seven months:

u	jan	feb	mar	apr	mai	jun	jul
$\text{pos}(u_1) + \text{pos}(u_2)$	11	11	14	17	14	31	31

Three colliding pairs already *before* the modulo: (jan, feb), (mar, mai), (jun, jul).