

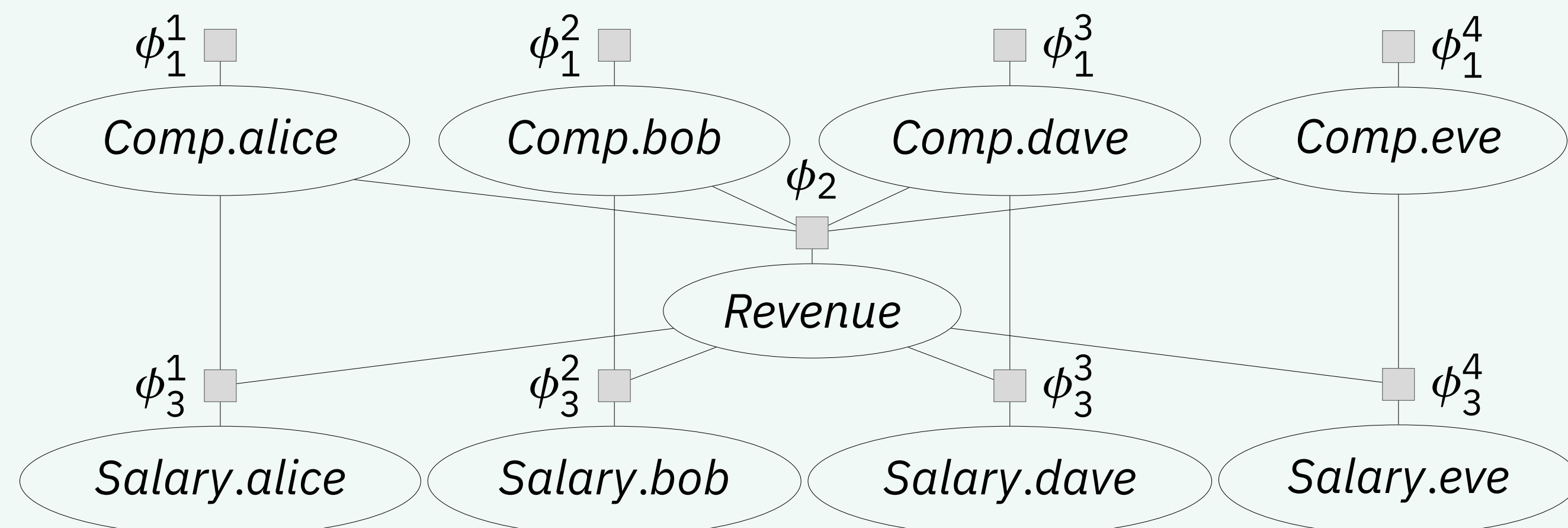
Efficient Detection of Commutative Factors in Factor Graphs

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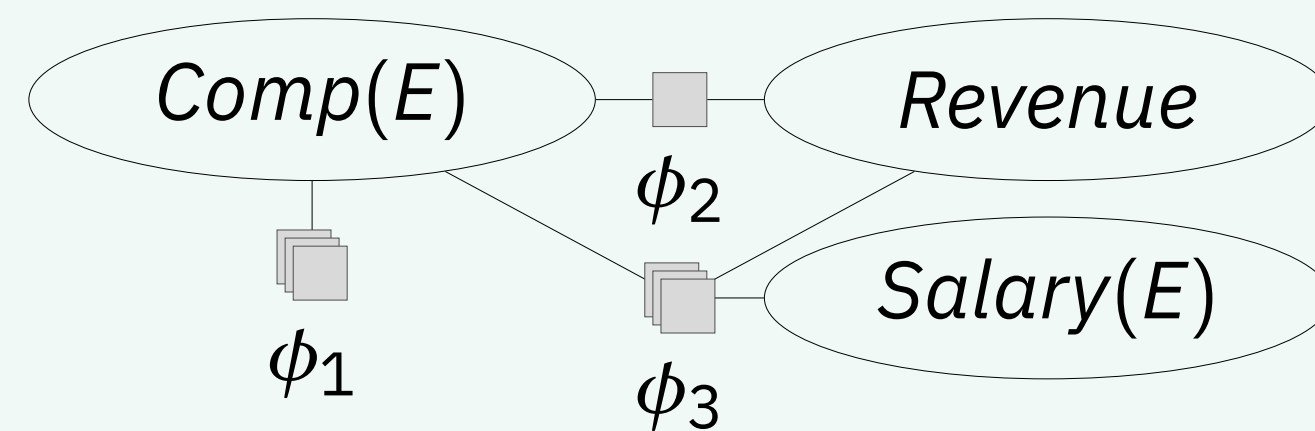
1. Motivation

- Factor graphs compactly encode a probability distribution
- Semantics of a factor graph G over a set of factors Φ :

$$P_G = \frac{1}{Z} \prod_{\phi \in \Phi} \phi$$



- Parametric factor graphs introduce logical variables to represent groups of random variables
- Parametric factor graphs enable lifted inference (idea: exploit indistinguishability of objects using exponentiation)



2. Problem Setup

- Goal: Efficiently detect commutative factors
- Commutative factors can be compressed

A	B	Rev	$\phi(A, B, Rev)$
true	true	true	φ_1
true	true	false	φ_4
true	false	true	φ_2
true	false	false	φ_5
false	true	true	φ_2
false	true	false	φ_5
false	false	true	φ_3
false	false	false	φ_6



$\#_E[Comp(E)]$	Rev	$\phi(\#_E[Comp(E)], Rev)$
[2, 0]	true	φ_1
[1, 1]	true	φ_2
[0, 2]	true	φ_3
[2, 0]	false	φ_4
[1, 1]	false	φ_5
[0, 2]	false	φ_6

3. Buckets

- Buckets count occurrences of range values in an assignment

A	B	R	$\phi(A, B, R)$	b
true	true	true	φ_1	[3, 0]
true	true	false	φ_4	[2, 1]
true	false	true	φ_2	[2, 1]
true	false	false	φ_5	[1, 2]
false	true	true	φ_2	[2, 1]
false	true	false	φ_5	[1, 2]
false	false	true	φ_3	[1, 2]
false	false	false	φ_6	[0, 3]

4. Detection of Commutative Factors (DECOR) Algorithm

Properties of buckets:

- For any subset S of commutative arguments:

$$|S| \leq \min_{b \in \{b | b \in \mathcal{B}(\phi) \wedge |\phi(b)| > 1\}} \max_{\varphi \in \phi(b)} \text{count}(\phi(b), \varphi)$$

(in each bucket b with $|\phi(b)| > 1$, there is a potential occurring at least $|S|$ times)

- For each group of identical potentials in a bucket, the intersection of their corresponding assignments yields candidates for commutative arguments—e.g., for φ_2 's in [2, 1]:

$$(true, false, true) \cap (false, true, true) = (\emptyset, \emptyset, true)$$

A	B	R	$\phi(A, B, R)$	b	$\phi(b)$
true	true	true	φ_1	[3, 0]	
true	true	false	φ_4	[2, 1]	
true	false	true	φ_2	[2, 1]	[3, 0] $\langle \varphi_1 \rangle$
true	false	false	φ_5	[1, 2]	[2, 1] $\langle \varphi_4, \varphi_2, \varphi_2 \rangle$
false	true	true	φ_2	[2, 1]	[1, 2] $\langle \varphi_5, \varphi_5, \varphi_3 \rangle$
false	true	false	φ_5	[1, 2]	[0, 3] $\langle \varphi_6 \rangle$
false	false	true	φ_3	[1, 2]	
false	false	false	φ_6	[0, 3]	

DECOR iterates over buckets and computes candidates for commutative arguments:

- Initial candidates: $\{\{A, B, R\}\}$
- Candidates for bucket [3, 0]: Skipped because $|\phi(b)| < 2$
- Candidates for bucket [2, 1]: $\{\{A, B\}\}$

$$(true, false, true) \cap (false, true, true) = (\emptyset, \emptyset, true)$$

- Candidates for bucket [1, 2]: $\{\{A, B\}\}$

$$(true, false, false) \cap (false, true, false) = (\emptyset, \emptyset, false)$$

- Candidates for bucket [0, 3]: Skipped because $|\phi(b)| < 2$
- Result (intersection over all buckets): $\{\{A, B\}\}$

5. Theoretical Guarantees

Given a factor $\phi(R_1, \dots, R_n)$:

- DECOR avoids »naive« iteration over all 2^n subsets of arguments
- Time complexity is upper-bounded depending on the number of groups of identical potentials in the buckets
- Expected worst-case time complexity of DECOR is in

$$O(|\mathcal{B}(\phi)| \cdot k \cdot \binom{n}{n/2} + r^n \cdot n),$$

where k is the maximum number of potentials within a bucket entailed by ϕ and r is the number of range values of $R_1, \dots, R_n$

6. Experiments

- Comparison of run times of DECOR and the »naive« approach
- Average run times over factors with $k \in \{0, 2, \lfloor \frac{n}{2} \rfloor, n-1, n\}$ commutative arguments
- Timeout after five minutes per instance

